



Using UniTime to build optimised academic schedules

Tomáš Müller June 2025



Agenda

What is UniTime

- A tool for building academic timetables/timetables
- Main component is a solver

Artificial Intelligence

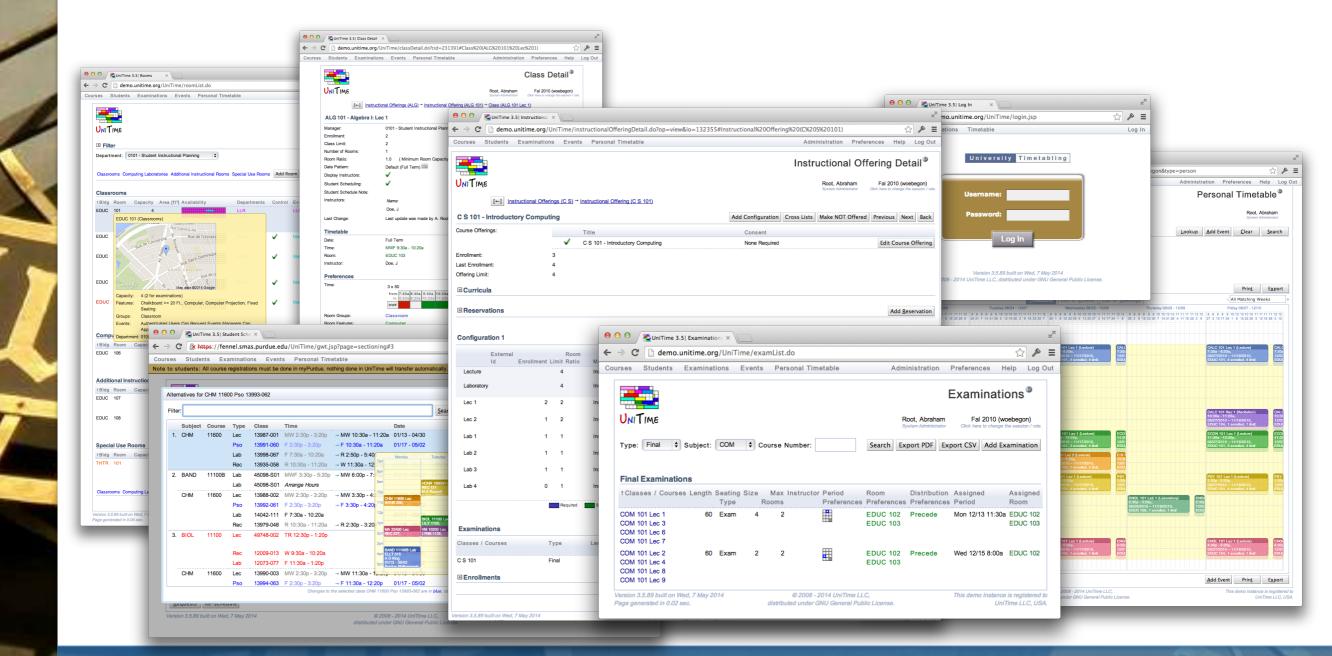
- Timetabling is a hard problem, with a lot of competing objectives
- Using a hybrid approach based on constraint programming (CP)
 - CP is a subfield of AI focusing on problem-solving

UniTime Solver

- Constraint Programming
- Algorithm basics
- A few details, pro & cons, etc.



• Comprehensive academic scheduling solution





- Comprehensive academic scheduling solution
- Five components

Courses: time & rooms

- time preferences
- room requirements
- availability
- course structure
- avoid time clashes
- student course demand
- distributions

Events: room bookings

- personal schedule
- approval process
- unavailabilities

- course-related events

Students: class schedule

Exams: period & rooms

period & room prefs

distributions

minimise student conflicts

room sharing or splitting

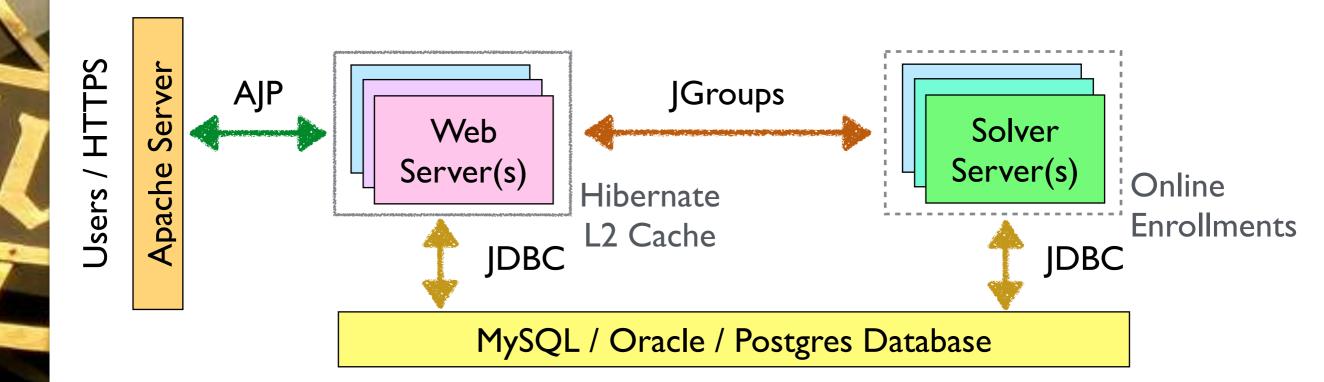
- course requirements
- alternatives / free times
- reservations
- section balancing
- schedule quality

Instructors: teaching assignments

- teaching load
- qualifications / skills
- preferences
- course spreading

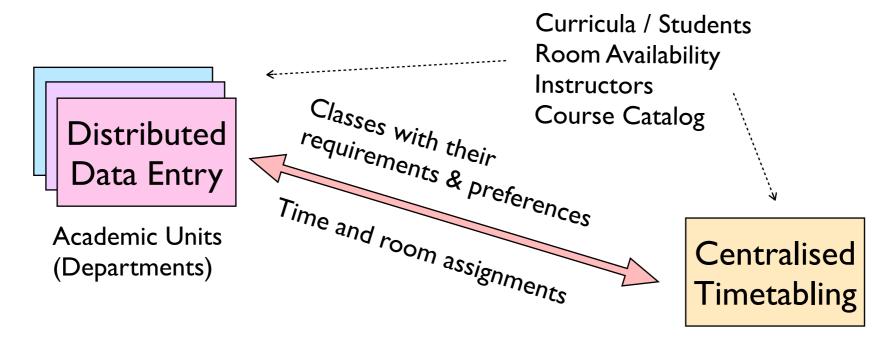


- Comprehensive academic scheduling solution
- Five components
- Open Source, web-based, written in Java using modern technologies





- Comprehensive academic scheduling solution
- Five components
- Open Source, web-based, written in Java using modern technologies
- Distributed data entry and timetabling in a multi-user environment



Timetabling Office (Registrar)



- Comprehensive academic scheduling solution
- Five components
- Open Source, web-based, written in Java using modern technologies
- Distributed data entry and timetabling in a multi-user environment
- Using state of the art optimisation algorithms

























ITC 2019

- course timetabling with students
- 3 sets of problems (10 each)
- 600+ registrations
- 20+ solvers
- 20+ publications

https://www.itc2019.org



Constraint Programming

Constraint Satisfaction Problem $\Theta = (V, D, C)$

- $V = \{v_1, v_2, ..., v_n\}$ is a finite set of variables
- D= $\{Dv_1, Dv_2, ..., Dv_n\}$ is a set of domains
 - Domain is a finite set of values
- C = $\{c_1, c_2, ..., c_m\}$ is a set of constraints
 - A constraint limits the combination of values that can variables simultaneously take
- Solution is an assignment of all variables η : $V \rightarrow D$
 - That satisfy all the constraints from C

Optimisation Problem $\Theta' = (V, D, C, f)$

- f is an objective function
 - That maps every partial feasible assignment to a number
 - Usually expressed by soft constraints and penalties



Examinations Example

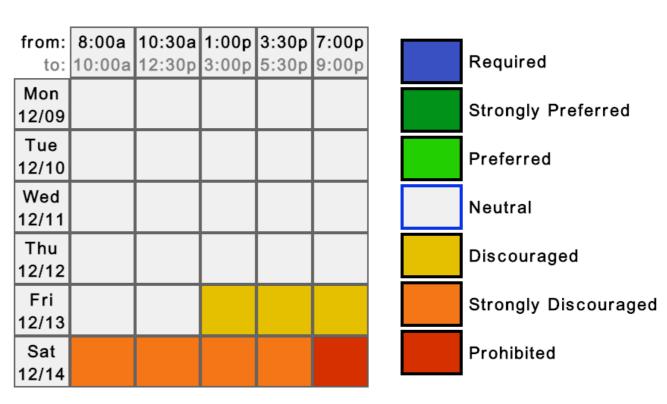
Example: Examination Timetabling

- Variables: individual examinations
- Domains: period and room for each exam
- Hard Constraints:
 - Examination size ≤ room capacity
 - Period and room requirements
 - No two exams in the same room at the same period
 - No student taking two exams during the same period

•

- Soft constraints:
 - Period and room preferences
 - Back-to-backs
 - Not more than 2 exams on a day
 - Large exams first

•





Timetabling is hard

Example: Examination Timetabling

- Variables: individual examinations
- Domains: period and room for each exam
- Hard Constraints:
 - Examination size ≤ room capacity
 - Period and room requirements
 - No two exams in the same room at the same period
 - No student taking two exams during the same period
 - •
- Soft constraints:
 - Period and room preferences
 - Back-to-backs
 - Not more than 2 exams on a day
 - Large exams first

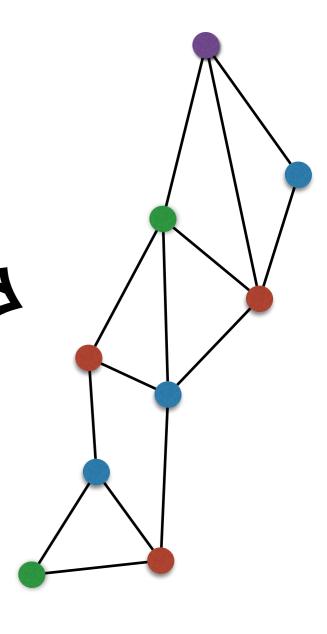
•

Graph colouring

Exams: vertices

Students: edges

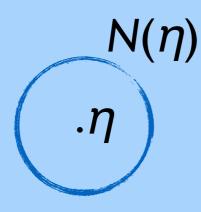
Colours: periods





Local Search

Search Space



In each iteration

- Generate a candidate solution $y \in N(x)$
- Evaluate a solution
 - Discard when conflicting
 - When accepted assign x = y
- Check if termination conditions are met

Neighbours

- Swap period/room
- Swap exams
- Assign exam, resolve/unassign conflicts
- Limited-depth search
- Kempe chain, MIP, ...

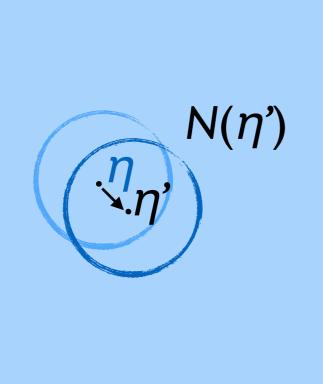
Acceptance criteria

- Hill Climbing $F(y) \le F(x)$
- Simulated Annealing $P(\text{accept}) = e^{-\Delta E/T}$
- Great Deluge $F(y) \leq Bound$



Local Search

Search Space



In each iteration

- Generate a candidate solution $y \in N(x)$
- Evaluate a solution
 - Discard when conflicting
 - When accepted assign x = y
- Check if termination conditions are met

Neighbours

- Swap period/room
- Swap exams
- Assign exam, resolve/unassign conflicts
- Limited-depth search
- Kempe chain, MIP, ...

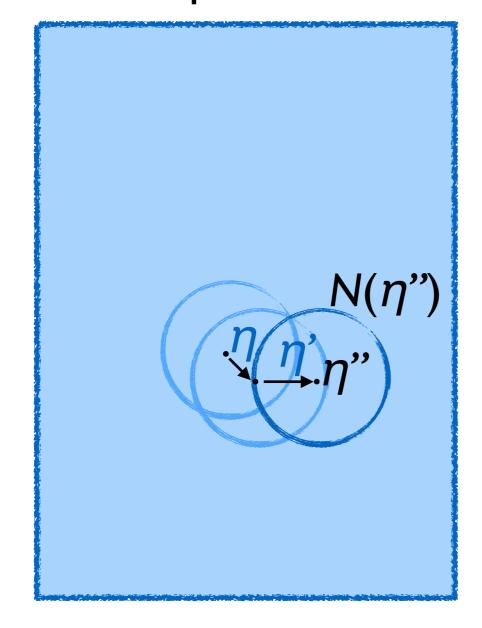
Acceptance criteria

- Hill Climbing $F(y) \le F(x)$
- Simulated Annealing $P(\text{accept}) = e^{-\Delta E/T}$
- Great Deluge $F(y) \leq Bound$



Local Search

Search Space



In each iteration

- Generate a candidate solution $y \in N(x)$
- Evaluate a solution
 - Discard when conflicting
 - When accepted assign x = y
- Check if termination conditions are met

Neighbours

- Swap period/room
- Swap exams
- Assign exam, resolve/unassign conflicts
- Limited-depth search
- Kempe chain, MIP, ...

Acceptance criteria

- Hill Climbing $F(y) \le F(x)$
- Simulated Annealing $P(\text{accept}) = e^{-\Delta E/T}$
- Great Deluge $F(y) \leq Bound$



Solver Framework



Course Timetabling

Examination Timetabling

Instructor Scheduling

Student Scheduling

Constraint-based Model

Variables

Values

Constraints

Criteria

Solver

Configuration

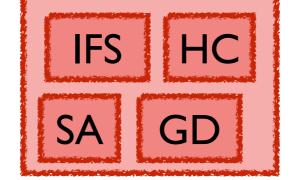
Parameters

General

Problem Specific

Neighbourhoods

Algorithms





UniTime Solver

Construction Phase

- Incomplete but feasible
- Pick an unassigned variable
- Find the best assignment
 - Unassign conflicts
- Conflict statistics
 - To prevent cycling
 - Provides feedback when over-constrained

Improvement Phase

- Complete and feasible
- Random selection of a candidate change
- Various acceptance criteria
 - HC, SA or GD
- Reheating
- Parameter adjustments
- Until time runs out

$$A \neq a \Leftarrow \begin{cases} 3 \times B = a \\ 4 \times B = c \\ 2 \times C = a \\ 120 \times D = a \end{cases}$$



Pros and Cons

Pros

- Usually very fast and scalable (compared to other methods)
- Flexible
 - Easy to extend
 - Exceptions are fairly easy to model
- Can start from an existing solution
 - Distance from previous solution can be also minimised
 - Can be used to provide suggestions (interactive timetabling)

Cons

- Cannot guarantee optimality
- Small changes can have large consequences
- Large changes may be expensive to make
- Often requires implementation of the model and at least a few problem specific heuristics/neighbourhoods



Other Approaches

Mixed Integer Programming (Operations Research)

- Using boolean or integer variables
- All constraints must be linear
- Existing (commercial) solvers that are getting very good
- Can guarantee optimality or provide upper bounds
- Heuristic approaches needed for large problems
 - Large neighbourhood search, fix & optimise, ruin & recreate

Comparison

- The first two places at ITC 2019 were using MIP solvers
 - Both using some version of fix & optimise
- Problem had to be segmented into smaller sub-problems
- The winner solver can provide better results than UniTime
 - But require lots of time and computational resources
- See ITC2019.org for more details



Conclusion

Thank you!

See www.unitime.org for more details about UniTime!

An online demo is available at https://demo.unitime.org